

CSE 331

Software Design and Implementation

Lecture 13

Generics<1>

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Varieties of abstraction

Abstraction over *computation*: procedures (methods)

```
int x1, y1, x2, y2;  
Math.sqrt(x1*x1 + y1*y1);  
Math.sqrt(x2*x2 + y2*y2);
```

Abstraction over *data*: ADTs (classes, interfaces)

```
Point p1, p2;
```

Abstraction over *types*: polymorphism (generics)

```
Point<Integer>, Point<Double>
```

Today!

Why we ❤ abstraction

Hide details

- Avoid distraction
- Permit details to change later

Give a *meaningful name* to a concept

Permit *reuse* in new contexts

- Avoid duplication: error-prone, confusing
- Save reimplementation effort
- Helps to “Don’t Repeat Yourself”



*Any true wizard knows, once you know
the name of a thing you can control it.*

-- Jerry Sussman

Related abstractions

```
interface ListOfStrings {  
    boolean add(String elt);  
    String get(int index);  
}  
  
interface ListOfNumbers {  
    boolean add(Number elt);  
    Number get(int index);  
}
```

Related abstractions

```
interface ListOfStrings {  
    boolean add(String elt);  
    String get(int index);  
}  
  
interface ListOfNumbers {  
    boolean add(Number elt);  
    Number get(int index);  
}
```

... and many, many more

```
// Type abstraction  
// abstracts over element type E  
interface List<E> {  
    boolean add(E n);  
    E get(int index);  
}
```

Type abstraction
lets us use these types:

```
List<String>  
List<Number>  
List<Integer>  
List<List<String>>  
...
```

Formal parameter vs. type parameter

```
interface ListOfIntegers {  
    boolean add(Integer elt);  
    Integer get(int index);  
}
```

- Declares a new **variable**, called a **(formal) parameter**
- **Instantiate** with any **expression** of the right type
 - E.g., `lst.add(7)`
- **Type** of `add` is `Integer → boolean`

```
interface List<E> {  
    boolean add(E n);  
    E get(int index);  
}
```

- Declares a new **type variable**, called a **type parameter**
- **Instantiate** with any (reference) type
 - E.g., `List<String>`
- **“Type”** of `List` is `Type → Type`
 - Never just use `List` (in Java for backward compatibility)

Type variables are types

```
Declaration  
class NewSet<T> implements Set<T> {  
    // rep invariant:  
    //     non-null, contains no duplicates  
    // ...  
    List<T> theRep;  
    T lastItemInserted;  
    ...  
}
```

The diagram illustrates the relationship between type variable declarations and their subsequent uses. It features two yellow rectangular boxes: one at the top labeled "Declaration" and one at the bottom labeled "Use". Four blue arrows originate from the declaration of each type variable and point to its corresponding use in the code. Specifically, the declaration of `T` at the top of the class body points to its use as the type of `lastItemInserted`. The declaration of `Set<T>` points to its use as the type parameter for `implements`. The declaration of `List<T>` points to its use as the type of `theRep`. The declaration of `theRep` points to its use as the type of `lastItemInserted`.

Declaring and instantiating generics

```
class MyClass<TypeVar1, ..., TypeVarN> {...}  
interface MyInterface<TypeVar1, ..., TypeVarN> {...}
```

- Convention: Type variable has one-letter name such as:
T for **Type**, **E** for **Element**,
K for **Key**, **V** for **Value**, ...

To instantiate a generic class/interface, client supplies *type arguments*:

```
MyClass<String, ..., Date>
```

Restricting instantiations by clients

```
boolean add1(Object elt);  
boolean add2(Number elt);  
add1(new Date()); // OK  
add2(new Date()); // compile-time error
```

Upper bounds

```
interface List1<E extends Object> {...}
```

```
interface List2<E extends Number> {...}
```

```
List1<Date> // OK, Date is a subtype of Object
```

```
List2<Date> // compile-time error, Date is not a  
// subtype of Number
```

Declaring and instantiating generics: syntax with bounds

```
class MyClass<TypeVar1 extends TypeBound1,  
             ...,  
             TypeVarN extends TypeBoundN> {...}
```

- (same for interface definitions)
- (default upper bound is `Object`)

To instantiate a generic class/interface, client supplies type arguments:

`MyClass<String, ..., Date>`

- Compile-time error if type is not a subtype of the upper bound

Using type variables

Code can perform any operation permitted by the bound

- Because we know all instantiations will be subtypes!
 - An enforced precondition on type instantiations

More examples

```
public class Graph<N> implements Iterable<N> {  
    private final Map<N, Set<N>> node2neighbors;  
    public Graph(Set<N> nodes, Set<Tuple<N,N>> edges) {  
        ...  
    }  
}  
  
public interface Path<N, P extends Path<N,P>>  
    extends Iterable<N>, Comparable<Path<?, ?>> {  
    public Iterator<N> iterator();  
    ...  
}
```

Do **NOT** copy/paste this stuff into your project unless it is what you want
– And you understand it!

More bounds

`<TypeVar extends SuperType>`

- An *upper bound*; accepts given supertype or any of its subtypes

`<TypeVar extends ClassA & InterfaceB & InterfaceC & ...>`

- *Multiple* upper bounds (superclass/interfaces) with &

Recursively-defined bounds:

```
// TreeSet works for any type that can be compared
// to itself.
```

```
public class TreeSet<T extends Comparable<T>> {  
    ...  
}
```

Outline

- Basics of generic types for classes and interfaces
- Basics of *bounding* generics
-  Generic *methods* [not just using type parameters of class]
- Generics and *subtyping*
- Using *bounds* for more flexible subtyping
- Using *wildcards* for more convenient bounds
- Related digression: Java's *array subtyping*
- Java realities: type erasure
 - Unchecked casts
 - `equals` interactions
 - Creating generic arrays

Generic classes are not enough

```
class Utils {  
    static double sumList(List<Number> lst) {  
        double result = 0.0;  
        for (Number n : lst) {  
            result += n.doubleValue();  
        }  
        return result;  
    }  
    static Object choose(List<Object> lst) {  
        int i = ... // random number < lst.size  
        return lst.get(i);  
    }  
}
```

Reminder: `static` means “no receiver (`this` parameter)”.

Cannot pass
`List<Double>`

Independent of
`Number` above

`List<Double>`
is not a subtype of
`List<Number>` !
We will see why soon.

Weaknesses of generic classes

- Would like to use `sumList` for any subtype of `Number`
 - For example, `Double` or `Integer`
 - But as we will see, `List<Double>` is not a subtype of `List<Number>`
- Would like to use `choose` for any element type
 - i.e., any subclass of `Object`
 - No need to restrict to subclasses of `Number`
 - Want to tell clients more about return type than `Object`
- Class `Utils` is not generic, but the *methods* should be generic

Generic methods solve the problem

```
class Utils {  
    static <T1 extends Number>  
        double sumList(List<T1> lst) {  
            double result = 0.0;  
            for (Number n : lst) { // T1 also works  
                result += n.doubleValue();  
            }  
            return result;  
        }  
    static <T2>  
        T2 choose(List<T2> lst) {  
            int i = ... // random number < lst.size  
            return lst.get(i);  
        }  
}
```

Have to declare type parameter(s)

Have to declare type parameter(s)

Using generics in methods

- Instance methods can use type parameters of the class
- Instance methods and static methods can have their own type parameters
 - Generic methods
- Callers to generic methods need not explicitly instantiate the methods' type parameters
 - Compiler usually figures it out for you
 - *Type inference*

More examples

```
<T extends Comparable<T>> T max(Collection<T> c) {  
    ...  
}
```

```
<T extends Comparable<T>>  
void sort(List<T> list) {  
    // ... use list.get() and T's compareTo  
}
```

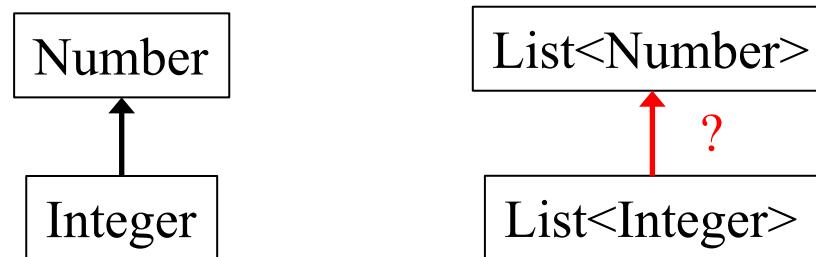
(This one works, but we will make it even more useful later by adding more bounds.)

```
<T> void copyTo(List<T> dst, List<T> src) {  
    for (T t : src)  
        dst.add(t);  
}
```

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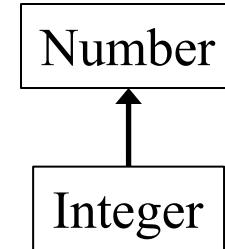
Generics and subtyping



- **Integer** is a subtype of **Number**
- Is **List<Integer>** a subtype of **List<Number>?**
- Use subtyping rules (stronger, weaker) to find out...

List<Number> and List<Integer>

```
interface List<T> {  
    boolean add(T elt);  
    T get(int index);  
}
```

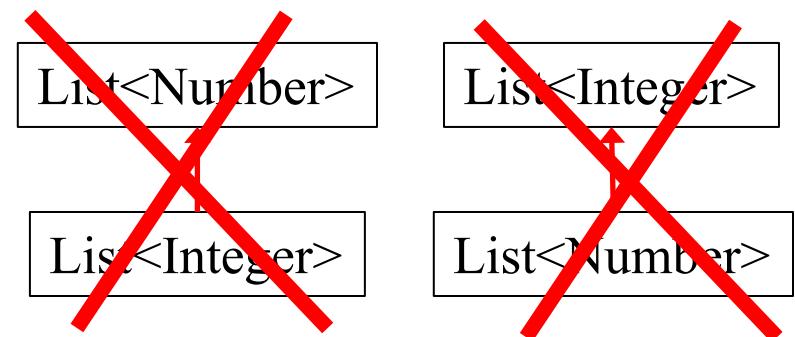


So type List<Number> has:

```
boolean add(Number elt);  
Number get(int index);
```

So type List<Integer> has:

```
boolean add(Integer elt);  
Integer get(int index);
```



Java subtyping is *invariant* with respect to generics

- Neither List<Number> nor List<Integer> subtype of other
- Not covariant and not contravariant

How to remember the invariant rule

If **Type2** and **Type3** are different,
then **Type1<Type2>** is *not* a subtype of **Type1<Type3>**

Previous example shows why:

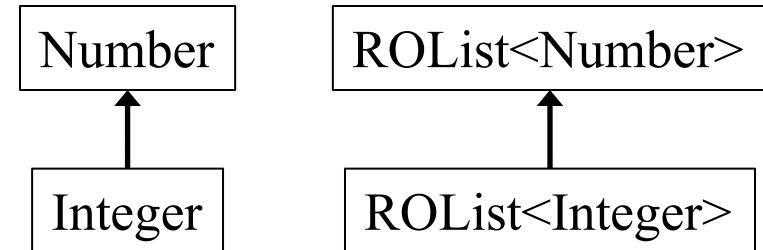
- Observer method prevents one direction
- Mutator/producer method prevents the other direction

If our types have only observers or only mutators, then one direction of subtyping would be sound

- Java's type system is not expressive enough to allow this

Read-only allows covariance

```
interface ReadOnlyList<T> {  
    T get(int index);  
}
```



Type `ReadOnlyList<Number>` has method:

```
Number get(int index);
```

Type `ReadOnlyList<Integer>` has method:

```
Integer get(int index);
```

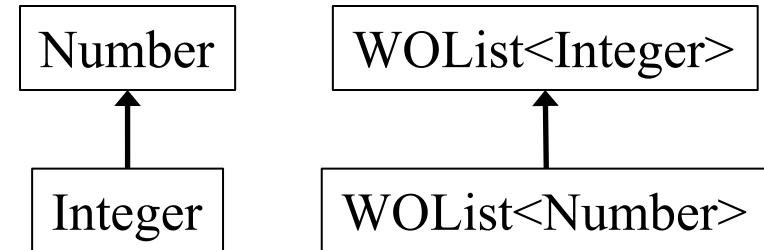
So *covariant* subtyping would be correct:

- `ROList<Integer>` is a subtype of `ROList<Number>`
- Covariant = type of `ROList<T>` changes the **same way** `T` changes

The Java type system conservatively disallows this subtyping

Write-only allows contravariance

```
interface WriteOnlyList<T> {  
    boolean add(T elt);  
}
```



Type `WriteOnlyList<Number>` has method:

```
boolean add(Number elt);
```

Type `WriteOnlyList<Integer>` has method:

```
boolean add(Integer elt);
```

So *contravariant* subtyping would be correct:

- `WOList<Number>` is a subtype of `WOList<Integer>`
- **Contra**variant = type of `WOList<T>` changes **opposite** to T

The Java type system conservatively disallows this subtyping

Generic types and subtyping

- `List<Integer>` and `List<Number>` are not subtype-related
- Generic types can have subtyping relationships
- Example: If `HeftyBag` extends `Bag`, then
 - `HeftyBag<Integer>` is a subtype of `Bag<Integer>`
 - `HeftyBag<Number>` is a subtype of `Bag<Number>`
 - `HeftyBag<String>` is a subtype of `Bag<String>`
 - ...

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