ADTs are defined by a specification

Abstract state + collection of procedural *abstractions*
  - Not a collection of procedures

Together, these procedural abstractions provide some *set of values*
  *All* the ways of directly using that set of values
  - Creating
  - Manipulating
  - Observing

- Creators and producers: make new values
- Mutators: change the value (but don't affect `==`)
- Observers: allow one to distinguish different values

ADTs and specifications

So far, we have only specified ADTs
  - Specification makes no reference to the implementation

Of course, we need [*guidelines for how*] to implement ADTs

Of course, we need [*guidelines for how*] to ensure our implementations satisfy our specifications

Two intellectual tools are really helpful…
Connecting implementations to specs

**Representation Invariant**: maps Object $\rightarrow$ boolean
- Indicates if an instance is *well-formed*
- Defines the set of valid concrete values
- Only values in the valid set make sense as implementations of an abstract value
- **For implementors/debuggers/maintainers of the abstraction**: no object should ever violate the rep invariant
  - Such an object has no useful meaning

**Abstraction Function**: maps Object $\rightarrow$ abstract value
- What the data structure *means* as an abstract value
- How the data structure is to be interpreted
- Only defined on objects meeting the rep invariant
- **For implementors/debuggers/maintainers of the abstraction**: Each procedure should meet its spec (abstract values) by "doing the right thing" with the concrete representation

Example: CharSet Abstraction

```java
// Overview: A CharSet is a finite mutable set of Characters
public CharSet() {...}
// @effects: creates a fresh, empty CharSet
public void insert(Character c) {...}
// @modifies: this
// @effects: this_post = this_pre + {c}
public void delete(Character c) {...}
// @modifies: this
// @effects: this_post = this_pre - {c}
public boolean member(Character c) {...}
// @return: cardinality of this
public int size() {...}
```

Implementing a Data Abstraction (ADT)

To implement a data abstraction:
- Select the representation of instances, “the rep”
  - In Java, typically instances of some class you define
- Implement operations in terms of that rep

Choose a representation so that:
- It is possible to implement required operations
- The most frequently used operations are efficient
  - But which will these be?
  - Abstraction allows the rep to change later

An implementation: Is it right?

```java
class CharSet {
private List<Character> elts =
new ArrayList<Character>();
public void insert(Character c) {
etls.add(c);
}
public void delete(Character c) {
etls.remove(c);
}
public boolean member(Character c) {
return elts.contains(c);
}
public int size() {
return elts.size();
}
}
CharSet s = new CharSet();
Character a = new Character('a');
s.insert(a);
s.insert(a);
s.delete(a);
if (s.member(a)) System.out.print("wrong");
else System.out.print("right");
Where is the error?
```
Where Is the Error?

If you can answer this, then you know what to fix

*Perhaps delete* is wrong
  – Should remove all occurrences?

*Perhaps insert* is wrong
  – Should not insert a character that is already there?

How can we know?
  – The *representation invariant* tells us
  – If it’s “our code”, this is how we document our choice for “the right answer”

Now we can locate the error

```java
public void insert(Character c) {
    elts.add(c);
}
public void delete(Character c) {
    elts.remove(c);
}
```

The representation invariant

- Defines data structure well-formedness
- Must hold before and after every CharSet operation
- Operations (methods) may depend on it
- Write it like this:

```java
class CharSet {
    // Rep invariant:
    //   elts has no nulls and no duplicates
    private List<Character> elts = ...
    ...

    // Or, more formally (if you prefer):
    ∀ indices i of elts . elts.elementAt(i) ≠ null
    ∀ indices i, j of elts .
    elts.elementAt(i).equals(elts.elementAt(j)) ⇒ i = j
```

Another example

```java
class Account {
    private int balance;
    // history of all transactions
    private List<Transaction> transactions;
    ...
}
```

Real-world constraints:
- Balance ≥ 0
- Balance = Σi transactions.get(i).amount

Implementation-related constraints:
- Transactions ≠ null
- No nulls in transactions
Checking rep invariants

Should code check that the rep invariant holds?

– Yes, if it’s inexpensive [depends on the invariant]
– Yes, for debugging [even when it’s expensive]
– Often hard to justify turning the checking off
– Some private methods need not check (Why?)

A great debugging technique:

*Design your code to catch bugs by implementing and using rep-invariant checking*

Practice defensive programming

Assume that you will make mistakes

Write and incorporate code designed to catch them

– On entry:
  • Check rep invariant
  • Check preconditions
– On exit:
  • Check rep invariant
  • Check postconditions

Checking the rep invariant helps you *discover* errors

Reasoning about the rep invariant helps you *avoid* errors

Listing the elements of a CharSet

Consider adding the following method to CharSet

// returns: a List containing the members of this
public List<Character> getElts();

Consider this implementation:

// Rep invariant: elts has no nulls and no dups
public List<Character> getElts() { return elts; }

Does the implementation of *getElts* preserve the rep invariant?

Kind of, sort of, not really....
Representation exposure

Consider this client code (outside the CharSet implementation):

```java
CharSet s = new CharSet();
Character a = new Character('a');
s.insert(a);
s.getElts().add(a);
s.delete(a);
if (s.member(a)) …
```

Representation exposure is external access to the rep

Representation exposure is almost always **EVIL**

*A BIG DEAL, A COMMON BUG, YOU NOW HAVE A NAME FOR IT!*

If you do it, document why and how

– And feel guilty about it!

Avoiding representation exposure

*Understand* what representation exposure is

*Design* ADT implementations to make sure it doesn’t happen

Treat rep exposure as a bug: *fix* your bugs

Test for it with *adversarial clients*:

– Pass values to methods and then mutate them
– Mutate values returned from methods

private is not enough

Making fields *private* does *not* suffice to prevent rep exposure

– See our example
– Issue is *aliasing of mutable data inside and outside the abstraction*

Examples of copying (assume Point is a mutable ADT):

```java
class Line {
    private Point s, e;
    public Line(Point s, Point e) {
        this.s = new Point(s.x,s.y);
        this.e = new Point(e.x,e.y);
    }
    public Point getStart() {
        return new Point(this.s.x,this.s.y);
    }
}
```

Avoiding rep exposure (way #1)

One way to avoid rep exposure is to make *copies* of all data that cross the abstraction barrier

– Copy in [parameters that become part of the implementation]
– Copy out [results that are part of the implementation]

*Two general ways to avoid representation exposure…*
Need deep copying

“Shallow” copying is not enough
– Prevent any aliasing to mutable data inside/outside abstraction

What's the bug (assuming Point is a mutable ADT)?
class PointSet {
  private List<Point> points = ...
  public List<Point> getElts() {
    return new ArrayList<Point>(points);
  }
}

Not in example: Also need deep copying on “copy in”

Why [not] immutability?

Several advantages of immutability
– Aliasing does not matter
– No need to make copies with identical contents
– Rep invariants cannot be broken
– See CSE341 for more!

Does require different designs (e.g., if Point immutable)
  void raiseLine(double deltaY) {
    this.s = new Point(s.x, s.y+deltaY);
    this.e = new Point(e.x, e.y+deltaY);
  }

Immutable classes in Java libraries include String, Character, Integer, ...

Avoiding rep exposure (way #2)

One way to avoid rep exposure is to exploit the immutability of (other) ADTs the implementation uses
– Aliasing is no problem if nobody can change data
  • Have to mutate the rep to break the rep invariant

Examples (assuming Point is an immutable ADT):
class Line {
  private Point s, e;
  public Line(Point s, Point e) {
    this.s = s;
    this.e = e;
  }
  public Point getStart() {
    return this.s;
  }
  ...
}

Deepness, redux

An immutable ADT must be immutable “all the way down”
– No references reachable to data that may be mutated

So combining our two ways to avoid rep exposure:
– Must copy-in, copy-out “all the way down” to immutable parts
Recall our initial rep-exposure example:

```java
class CharSet {
    // Rep invariant: elts has no nulls and no dups
    private List<Character> elts = ...;

    // returns: elts currently in the set
    public List<Character> getElts() {
        return new ArrayList<Character>(elts); //copy out!
    }
    ...
}
```

An alternative

```java
// returns: elts currently in the set
public List<Character> getElts() { // version 1
    return new ArrayList<Character>(elts); //copy out!
}
public List<Character> getElts() { // version 2
    return Collections.unmodifiableList<Character>(elts);
}
```

From the JavaDoc for `Collections.unmodifiableList`:

> Returns an unmodifiable view of the specified list. This method allows modules to provide users with "read-only" access to internal lists. Query operations on the returned list "read through" to the specified list, and attempts to modify the returned list... result in an UnsupportedOperationException.

The good news

```java
public List<Character> getElts() { // version 2
    return Collections.unmodifiableList<Character>(elts);
}
```

Clients cannot modify (mutate) the rep

- So they cannot break the rep invariant

(For long lists,) more efficient than copy out

Uses standard libraries

The bad news

```java
public List<Character> getElts() { // version 1
    return new ArrayList<Character>(elts); //copy out!
}
public List<Character> getElts() { // version 2
    return Collections.unmodifiableList<Character>(elts);
}
```

The two implementations do not do the same thing!

- Both avoid allowing clients to break the rep invariant
- Both return a list containing the elements

But consider:

```java
xs = s.getElts();
s.insert('a');
xs.contains('a');
```

Version 2 is observing an exposed rep, leading to different behavior.
Different specifications

Ambiguity of “returns a list containing the current set elements”

“returns a fresh mutable list containing the elements in the set

at the time of the call”

versus

“returns read-only access to a list that the ADT

continues to update to hold the current elements in the set”

A third spec weaker than both [but less simple and useful!]

“returns a list containing the current set elements. Behavior is

unspecified (!) if client attempts to mutate the list or to access the list

after the set’s elements are changed”

Also note: Version 2’s spec also makes changing the rep later harder

– Only “simple” to implement with rep as a List