

Pumping Lemma

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May 17

Announcements

- Handouts
 - Feedback form
 - Last "unofficial" feedback
 - Solutions to H/W #5
- Office hours next week
 - Monday, May 22nd : no office hours
 - Wednesday, May 24th : 3:40-4:30pm
- No puzzle today

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1 (d) on H/W #6

- Should read exactly one extra (
 - Problem say one extra)

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Jenny's office hours

- No office hour today
- Tomorrow, Thu May 18, 1:00 – 3:00pm
- 4th floor breakout area

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Last lecture

- Converting PDAs to CFG
- Given $M = \langle Q, \Sigma, \Gamma, \delta, s, f \rangle$
- Build a $G = \langle V, \Sigma, R, A_{sf} \rangle$
 - $V = \{A_{pq} \mid p, q \in Q\}$
- Three kinds of rules
 - For every $p, q, r, s \in Q, t \in \Gamma, a, b \in \Sigma \cup \{\epsilon\}$, if 

```
graph LR; p((p)) -- "a, ε → t" --> r((r)); s((s)) -- "b, t → ε" --> q((q))
```
 - then add the rule $A_{pq} \rightarrow aA_{rs}b$

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The other rules

- For every $p, q, r \in Q$
 - $A_{pq} \rightarrow A_{pr}A_{rq}$
- For every $p \in Q$
 - $A_{pp} \rightarrow \epsilon$
- Recall $A_{pq} \Rightarrow^* w$ "signifies" the following:



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Idea (hope ?) behind the rules

- $A_{pq} \rightarrow aA_{rs}b$



- $A_{pq} \rightarrow A_{pr}A_{rq}$



Why does this work ?



We need a proof...

- Today briefly sketch one direction
- If A_{pq} generates a string then



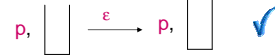
- See Claim 2.30, 2.31 in Sipser (Pgs 121-122)
 - Formal proof by induction for both directions

$$A_{pq} \Rightarrow^* w$$

- Induction on number of steps in the derivation

- Base case: 1 step in the derivation

- Only possible rule is $A_{pp} \rightarrow \epsilon$



The induction hypothesis

- For any $p, q \in Q$, if $A_{pq} \Rightarrow^* w$ in $\leq k$ steps then



- Inductive step: $A_{pq} \Rightarrow^* w$ in $k+1$ steps

- The derivation has to look like either
- $A_{pq} \rightarrow aA_{rs}b \Rightarrow^* ayb = w$, or
- $A_{pq} \rightarrow A_{pr}A_{rq} \Rightarrow^* yz = w$

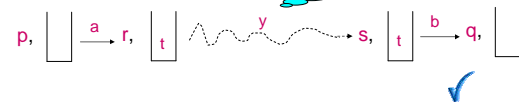
Case 1: $A_{pq} \rightarrow aA_{rs}b \Rightarrow^* ayb = w$

- $A_{rs} \Rightarrow^* y$ takes k steps

- By induction hypothesis



- This also implies



Case 2: $A_{pq} \rightarrow A_{pr}A_{rq} \Rightarrow^* yz = w$

■ By induction



■ Thus,



Questions ?

Up next: Ms. PDA finds out her limits



Screenshot by Kirk Geiger @kgeiger.net

And we will see our friend...

■ (Another) pumping lemma

