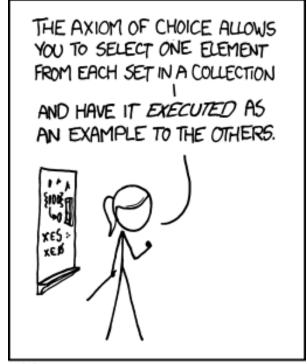
# **CSE 311: Foundations of Computing**

### **Lecture 8: Predicate Logic Proofs**



MY MATH TEACHER WAS A BIG BELIEVER IN PROOF BY INTIMIDATION.

## **Last class: Propositional Inference Rules**

Two inference rules per binary connective, one to eliminate it and one to introduce it

Elim ∧ 
$$A \land B$$
  
∴ A, B

Intro ∧  $A; B$   
∴ A ∧ B

Elim ∨  $A \lor B; \neg A$   
∴ B

Intro ∨  $A \lor B, B \lor A$ 

Modus Ponens  $A; A \to B$   
∴ B

Direct Proof  $A \Rightarrow B$   
Rule

∴  $A \to B$ 

Not like other rules

### Last class: Example

Prove: 
$$((p \rightarrow q) \land (q \rightarrow r)) \rightarrow (p \rightarrow r)$$

1.1. 
$$(p \rightarrow q) \land (q \rightarrow r)$$
 Assumption  
1.2.  $p \rightarrow q$   $\land$  Elim: 1.1  
1.3.  $q \rightarrow r$   $\land$  Elim: 1.1  
1.4.1.  $p$  Assumption  
1.4.2.  $q$  MP: 1.2, 1.4.1  
1.4.3.  $r$  MP: 1.3, 1.4.2  
1.4.  $p \rightarrow r$  Direct Proof Rule  
1.  $((p \rightarrow q) \land (q \rightarrow r)) \rightarrow (p \rightarrow r)$  Direct Proof Rule

# Last class: One General Proof Strategy

- 1. Look at the rules for introducing connectives to see how you would build up the formula you want to prove from pieces of what is given
- 2. Use the rules for eliminating connectives to break down the given formulas so that you get the pieces you need to do 1.
- 3. Write the proof beginning with what you figured out for 2 followed by 1.

### **Last Class: Some Inference Rules for Quantifiers**

P(c) for some c 
$$\exists x P(x)$$
 $\exists x P(x)$ 
 $\exists x P(x)$ 
 $\exists x P(x)$ 

# **Last Class: Predicate Logic Proofs**

- Can use
  - Predicate logic inference rules whole formulas only
  - Predicate logic equivalences (De Morgan's)
     even on subformulas
  - Propositional logic inference rules whole formulas only
  - Propositional logic equivalences
     even on subformulas

# **Last Class: Predicate Logic Proof**

$$\begin{array}{c}
P(c) \text{ for some c} \\
\therefore \quad \exists x P(x)
\end{array}$$

 $\begin{array}{c}
\forall x \ P(x) \\
\therefore \ P(a) \ \text{for any } a
\end{array}$ 

Prove  $\forall x P(x) \rightarrow \exists x P(x)$ 

- **1.1.**  $\forall x P(x)$  Assumption
- 1.2 P(a) Elim  $\forall$ : 1.1
- **1.3.**  $\exists x P(x)$  Intro  $\exists : 1.2$
- 1.  $\forall x P(x) \rightarrow \exists x P(x)$  Direct Proof Rule

# Inference Rules for Quantifiers: First look

P(c) for some c
$$\therefore \exists x P(x)$$

$$\begin{array}{c}
\forall x P(x) \\
\therefore P(a) \text{ for any } a
\end{array}$$

Let a be arbitrary\*"...P(a)
∴ ∀x P(x)

\* in the domain of P

 $\exists x P(x)$ 

∴ P(c) for some special\*\* c

\*\* By special, we mean that c is a name for a value where P(c) is true. We can't use anything else about that value, so c has to be a NEW name!

## **Predicate Logic Proofs with more content**

- In propositional logic we could just write down other propositional logic statements as "givens"
- Here, we also want to be able to use domain knowledge so proofs are about something specific
- Example: Domain of Discourse Integers
- Given the basic properties of arithmetic on integers,
   define:

  Predicate Definitions

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
Odd(x)  $\equiv \exists y (x = 2 \cdot y + 1)$ 

# A Not so Odd Example

### Domain of Discourse

Integers

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
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Prove "There is an even number"

Formally: prove  $\exists x \; Even(x)$ 

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Prove "There is an even number"

Formally: prove  $\exists x \; Even(x)$ 

1. 
$$2 = 2 \cdot 1$$
 Arithmetic

**2.** 
$$\exists y (2 = 2 \cdot y)$$
 Intro  $\exists : 1$ 

4. 
$$\exists x \; Even(x)$$
 Intro  $\exists : 3$ 

## A Prime Example

#### **Domain of Discourse**

Integers

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
Odd(x)  $\equiv \exists y (x = 2 \cdot y + 1)$   
Prime(x)  $\equiv "x > 1$  and  $x \ne a \cdot b$  for  
all integers a, b with  $1 < a < x$ "

**Prove "There is an even prime number"** 

### A Prime Example

#### **Domain of Discourse**

Integers

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all integers a, b with  $1 < a < x$ "

Prove "There is an even prime number"

Formally: prove  $\exists x (Even(x) \land Prime(x))$ 

- 1. 2 = 2.1
- **2.** Prime(**2**)\*

**Arithmetic** 

**Property of integers** 

<sup>\*</sup> Later we will further break down "Prime" using quantifiers to prove statements like this

### A Prime Example

#### **Domain of Discourse**

Integers

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$
  
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all integers a, b with  $1 < a < x$ "

Prove "There is an even prime number"

Formally: prove  $\exists x (Even(x) \land Prime(x))$ 

- 1.  $2 = 2 \cdot 1$  Arithmetic
- 2. Prime(2)\* Property of integers
- 3.  $\exists y (2 = 2 \cdot y)$  Intro  $\exists : 1$
- 4. Even(2) Defn of Even: 3
- 5. Even(2)  $\land$  Prime(2) Intro  $\land$ : 2, 4
- 6.  $\exists x (Even(x) \land Prime(x))$  Intro  $\exists : 5$

<sup>\*</sup> Later we will further break down "Prime" using quantifiers to prove statements like this

# Inference Rules for Quantifiers: First look

P(c) for some c
$$\therefore \exists x P(x)$$

$$\begin{array}{c}
\forall x P(x) \\
\therefore P(a) \text{ for any } a
\end{array}$$

Let a be arbitrary\*"...P(a)
∴ ∀x P(x)

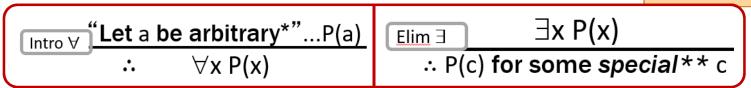
\* in the domain of P

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\*\* By special, we mean that c is a name for a value where P(c) is true. We can't use anything else about that value, so c has to be a NEW name!

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers

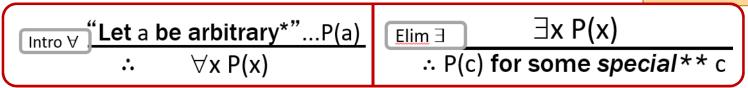


Prove: "The square of every even number is even."

Formal proof of:  $\forall x \text{ (Even}(x) \rightarrow \text{Even}(x^2))$ 

3.  $\forall x \text{ (Even}(x) \rightarrow \text{Even}(x^2))$ 

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers



Prove: "The square of every even number is even."

Formal proof of:  $\forall x \text{ (Even}(x) \rightarrow \text{Even}(x^2))$ 

1. Let a be an arbitrary integer

- 2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)
- 3.  $\forall x (Even(x) \rightarrow Even(x^2))$



Intro ∀: 1,2

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers



Prove: "The square of every even number is even."

Formal proof of:  $\forall x \text{ (Even}(x) \rightarrow \text{Even}(x^2))$ 

1. Let a be an arbitrary integer

**2.1** Even(a)

Assumption

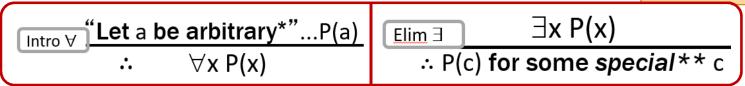
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Direct proof rule

Intro ∀: 1,2

Even(x)  $\equiv \exists y (x=2y)$  $Odd(x) \equiv \exists y (x=2y+1)$ Domain: Integers



Prove: "The square of every even number is even."

Formal proof of:  $\forall x (Even(x) \rightarrow Even(x^2))$ 

1. Let a be an arbitrary integer

2.2 
$$\exists y (a = 2y)$$
 Definition of Even

**2.5** 
$$\exists y (a^2 = 2y)$$

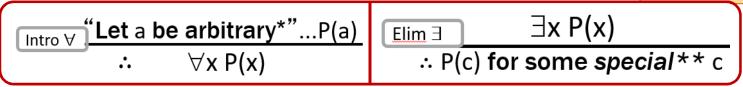
2. Even(a)
$$\rightarrow$$
Even(a<sup>2</sup>) Direct proof rule

3. 
$$\forall x (Even(x) \rightarrow Even(x^2))$$

**Definition of Even** 

Intro  $\forall$ : 1,2

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers



Prove: "The square of every even number is even."

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**Definition of Even** 

Need 
$$a^2 = 2c$$
 for some c

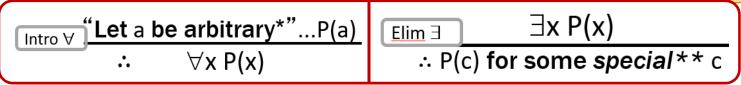
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Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers

Need  $a^2 = 2c$ 

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Prove: "The square of every even number is even."

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$$a = 2b$$
 Elim  $\exists$ : b special depends on a

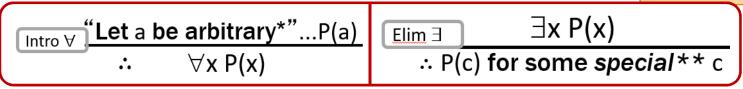
**2.5** 
$$\exists y (a^2 = 2y)$$
 Intro  $\exists rule: ?$ 

2. Even(a)
$$\rightarrow$$
Even(a<sup>2</sup>) Direct proof rule

3. 
$$\forall x (Even(x) \rightarrow Even(x^2))$$
 Intro  $\forall : 1,2$ 

Even(x)  $\equiv \exists y \ (x=2y)$ Odd(x)  $\equiv \exists y \ (x=2y+1)$ Domain: Integers

Used  $a^2 = 2c$  for  $c=2b^2$ 



Prove: "The square of every even number is even."

Formal proof of:  $\forall x \text{ (Even}(x) \rightarrow \text{Even}(x^2))$ 

1. Let a be an arbitrary integer

2.1	Even(a)	Assumption
2.2	$\exists y (a = 2y)$	Definition of Even

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**2.4** 
$$a^2 = 4b^2 = 2(2b^2)$$
 Algebra

**2.5** 
$$∃y (a^2 = 2y)$$
 Intro  $∃$  rule

2. Even(a)
$$\rightarrow$$
Even(a<sup>2</sup>) Direct proof rule

3. 
$$\forall x (Even(x) \rightarrow Even(x^2))$$
 Intro  $\forall : 1,2$ 

### Why did we need to say that b depends on a?

There are extra conditions on using these rules:

Let a be arbitrary\*"...P(a)

∴ 
$$\forall x P(x)$$

\* in the domain of P

Elim∃  $\exists x P(x)$ 

∴  $P(c)$  for some special\*\* c

\*\* c has to be a NEW name.

Over integer domain:  $\forall x \exists y (y \ge x)$  is True but  $\exists y \forall x (y \ge x)$  is False

#### **BAD "PROOF"**

- **1.**  $\forall x \exists y (y \ge x)$  Given
- 2. Let a be an arbitrary integer
- 3.  $\exists y (y \ge a)$  Elim  $\forall : 1$
- 4.  $b \ge a$  Elim  $\exists$ : b special depends on a
- 5.  $\forall x (b \ge x)$  Intro  $\forall : 2,4$
- 6.  $\exists y \forall x (y \ge x)$  Intro  $\exists : 5$

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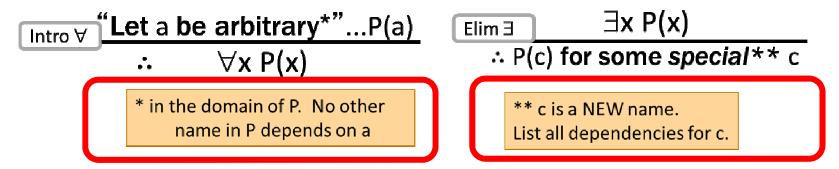
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Can't get rid of a since another name in the same line, b, depends on it!

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Can't get rid of a since another name in the same line, b, depends on it!

# Inference Rules for Quantifiers: Full version

P(c) for some c
$$\therefore \exists x P(x)$$

$$\begin{array}{c}
\forall x \ P(x) \\
\therefore \ P(a) \ \text{for any } a
\end{array}$$

Intro  $\forall$  "Let a be arbitrary\*"...P(a)
∴  $\forall x P(x)$ 

\* in the domain of P. No other name in P depends on a

 $\exists x P(x)$ 

∴ P(c) for some special\*\* c

\*\* c is a NEW name. List all dependencies for c.

# **English Proofs**

- We often write proofs in English rather than as fully formal proofs
  - They are more natural to read

- English proofs follow the structure of the corresponding formal proofs
  - Formal proof methods help to understand how proofs really work in English...
    - ... and give clues for how to produce them.

# **An English Proof**

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y (x = 2 \cdot y)$$

$$Odd(x) \equiv \exists y (x = 2 \cdot y + 1)$$

### **Prove "There is an even integer"**

#### **Proof:**

$$2 = 2 \cdot 1$$

1. 2 = 2.1

**Arithmetic** 

**2.**  $\exists y (2 = 2 \cdot y)$  Intro  $\exists : 1$ 

Therefore 2 is even.



3. Even(2)

Defn of Even: 2

Therefore, there is an even integer



 $\bullet$  4.  $\exists x \, Even(x)$ 

Intro ∃: 3

# **English Even and Odd**

Even(x)  $\equiv \exists y (x=2y)$  $Odd(x) \equiv \exists y (x=2y+1)$ Domain: Integers

Prove "The square of every even integer is even."

even integer.

Proof: Let a be an arbitrary 1. Let a be an arbitrary integer

2.1 Even(a) Assumption

Then, by definition, a = 2bfor some integer b (depending on a).

2.2 ∃y (a = 2y) Definition

2.3 a = 2b**b** special depends on **a** 

Squaring both sides, we get  $2.4 \text{ a}^2 = 4b^2 = 2(2b^2)$  Algebra  $a^2 = 4b^2 = 2(2b^2)$ .

Since 2b<sup>2</sup> is an integer, by definition, a<sup>2</sup> is even.

2.5  $\exists y (a^2 = 2y)$ 

2.6 Even(a<sup>2</sup>) Definition

Since a was arbitrary, it follows that the square of every even number is even. 2. Even(a) $\rightarrow$ Even(a<sup>2</sup>)

3.  $\forall x \text{ (Even(x)} \rightarrow \text{Even(x}^2\text{))}$ 

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y \ (x = 2y)$$
  
Odd(x)  $\equiv \exists y \ (x = 2y + 1)$ 

Domain of Discourse Integers

Prove "The square of every odd number is odd."

#### **Predicate Definitions**

Even(x) 
$$\equiv \exists y \ (x = 2y)$$
  
Odd(x)  $\equiv \exists y \ (x = 2y + 1)$ 

Domain of Discourse
Integers

# Prove "The square of every odd number is odd."

**Proof:** Let b be an arbitrary odd number.

Then, b = 2c+1 for some integer c (depending on b).

Therefore,  $b^2 = (2c+1)^2 = 4c^2 + 4c + 1 = 2(2c^2 + 2c) + 1$ .

Since  $2c^2+2c$  is an integer,  $b^2$  is odd. The statement follows since b was arbitrary.

### **Proofs**

- Formal proofs follow simple well-defined rules and should be easy to check
  - In the same way that code should be easy to execute
- English proofs correspond to those rules but are designed to be easier for humans to read
  - Easily checkable in principle