## CSE 311: Foundations of Computing I

## Section 8: CFGs, Relations, DFAs Solutions

## 1. CFGs

Construct CFGs for the following languages:
(a) All binary strings that end in 00 .

## Solution:

$$
\mathbf{S} \rightarrow 0 \mathbf{S}|1 \mathbf{S}| 00
$$

(b) All binary strings that contain at least three 1's.

## Solution:

$$
\begin{aligned}
& \mathbf{S} \rightarrow \mathbf{T} \mathbf{T} \\
& \mathbf{T} \rightarrow 0 \mathbf{T}|\mathbf{T} 0| 1 \mathbf{T} \mid 1
\end{aligned}
$$

(c) All binary strings with an equal number of 1 's and 0 's.

## Solution:

$$
\mathbf{S} \rightarrow 0 \mathbf{S} 1 \mathbf{S}|1 \mathbf{S} 0 \mathbf{S}| \varepsilon
$$

and

$$
\mathbf{S} \rightarrow \mathbf{S S}|0 \mathbf{S} 1| 1 \mathbf{S} 0 \mid \varepsilon
$$

both work. Note: The fact that all the strings generated have the property is easy to show (by induction) but the fact that one can generate all strings with the property is trickier. To argue this that each of these is grammars is enough one would need to consider how the difference between the \# of 0 's seen and the \# of 1 's seen occurs in prefixes of any string with the property.

## 2. Relations

(a) Draw the transitive-reflexive closure of $\{(1,2),(2,3),(3,4)\}$.

## Solution:


(b) Suppose that $R$ is reflexive. Prove that $R \subseteq R^{2}$.

## Solution:

Suppose $(a, b) \in R$. Since $R$ is reflexive, we know $(b, b) \in R$ as well. Since there is a $b$ such that $(a, b) \in R$ and $(b, b) \in R$, it follows that $(a, b) \in R^{2}$. Thus, $R \subseteq R^{2}$.
(c) Consider the relation $R=\{(x, y): x=y+1\}$ on $\mathbb{N}$. Is $R$ reflexive? Transitive? Symmetric? Anti-symmetric?

## Solution:

It isn't reflexive, because $1 \neq 1+1$; so, $(1,1) \notin R$. It isn't symmetric, because $(2,1) \in R$ (because $2=1+1$ ), but $(1,2) \notin R$, because $1 \neq 2+1$. It isn't transitive, because note that $(3,2) \in R$ and $(2,1) \in R$, but $(3,1) \notin R$. It is anti-symmetric, because consider $(x, y) \in R$ such that $x \neq y$. Then, $x=y+1$ by definition of $R$. However, $(y, x) \notin R$, because $y=x-1 \neq x+1$.
(d) Consider the relation $S=\left\{(x, y): x^{2}=y^{2}\right\}$ on $\mathbb{R}$. Prove that $S$ is reflexive, transitive, and symmetric.

## Solution:

Consider $x \in \mathbb{R}$. Note that by definition of equality, $x^{2}=x^{2}$; so, $(x, x) \in R$; so, $R$ is reflexive.

Consider $(x, y) \in R$. Then, $x^{2}=y^{2}$. It follows that $y^{2}=x^{2}$; so, $(y, x) \in R$. So, $R$ is symmetric.
Suppose $(x, y) \in R$ and $(y, z) \in R$. Then, $x^{2}=y^{2}$, and $y^{2}=z^{2}$. Since equality is transitive, $x^{2}=z^{2}$. So, $(x, z) \in R$. So, $R$ is transitive.

## 3. DFAs, Stage 1

Construct DFAs to recognize each of the following languages. Let $\Sigma=\{0,1,2,3\}$.
(a) All binary strings.

## Solution:


$q_{0}$ : binary strings
$q_{1}$ : strings that contain a character which is not 0 or 1 .
(b) All strings whose digits sum to an even number.

## Solution:


(c) All strings whose digits sum to an odd number.

## Solution:



## 4. DFAs, Stage 2

Construct DFAs to recognize each of the following languages. Let $\Sigma=\{0,1\}$.
(a) All strings which do not contain the substring 101.

## Solution:


$q_{3}$ : string that contain 101.
$q_{2}$ : strings that don't contain 101 and end in 10.
$q_{1}$ : strings that don't contain 101 and end in 1.
$q_{0}: \varepsilon, 0$, strings that don't contain 101 and end in 00.
(b) All strings containing at least two 0 's and at most one 1 .

## Solution:


(c) All strings containing an even number of 1 's and an odd number of 0 's and not containing the substring 10.

## Solution:



