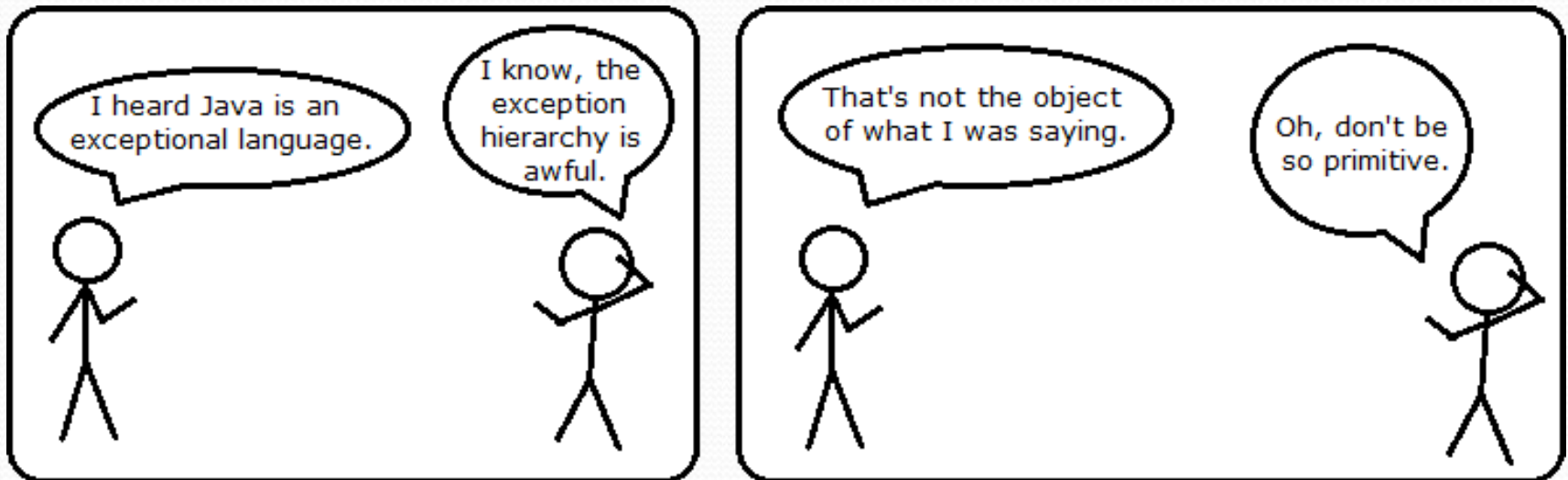


CSE 143

Lecture 4: testing and complexity

reading: 13.1-13.2



Tips for testing

- You cannot test every possible input, parameter value, etc.
 - Think of a limited set of tests likely to expose bugs.
- Think about boundary cases
 - Positive; zero; negative numbers
 - Right at the edge of an array or collection's size
- Think about empty cases and error cases
 - 0, -1, null; an empty list or array
- test behavior in combination
 - Maybe `add` usually works, but fails after you call `remove`
 - Make multiple calls; maybe `size` fails the second time only

Interfaces

- **interface:** A list of methods that a class can promise to implement.
 - Inheritance gives you an is-a relationship *and* code sharing.
 - A `Lawyer` can be treated as an `Employee` and inherits its code.
 - Interfaces give you an is-a relationship *without* code sharing.
 - A `Rectangle` object can be treated as a `Shape` but inherits no code.
 - Always declare variables using the ***interface*** type.

```
List<String> list = new ArrayList<String>();
```

Runtime Efficiency (13.2)

- **efficiency:** measure of computing resources used by code.
 - can be relative to speed (time), memory (space), etc.
 - most commonly refers to run time
- Assume the following:
 - Any single Java statement takes same amount of time to run.
 - A method call's runtime is measured by the total of the statements inside the method's body.
 - A loop's runtime, if the loop repeats N times, is N times the runtime of the statements in its body.

Efficiency examples

```
statement1;  
statement2;  
statement3;
```

} 3

```
for (int i = 1; i <= N; i++) {  
    statement4;  
}
```

} N

```
for (int i = 1; i <= N; i++) {  
    statement5;  
    statement6;  
    statement7;  
}
```

} 3N

} 4N + 3

Efficiency examples 2

```
for (int i = 1; i <= N; i++) {  
    for (int j = 1; j <= N; j++) {  
        statement1;  
    }  
}
```

} N^2

```
for (int i = 1; i <= N; i++) {  
    statement2;  
    statement3;  
    statement4;  
    statement5;  
}
```

} $4N$

} $N^2 + 4N$

- How many statements will execute if $N = 10$? If $N = 1000$?

Algorithm growth rates (13.2)

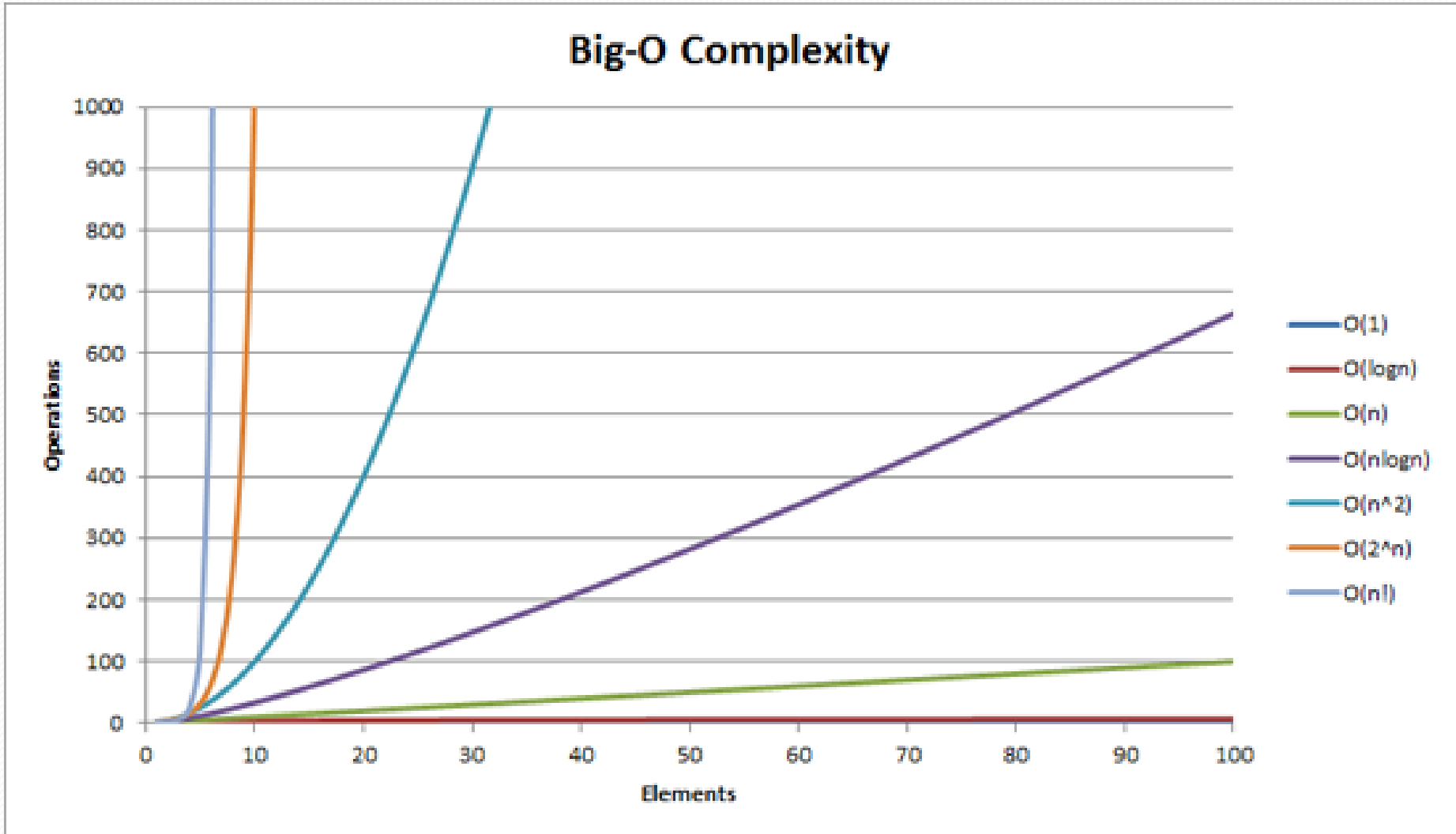
- We measure runtime in proportion to the input data size, N .
 - **growth rate**: Change in runtime as N changes.
- Say an algorithm runs $0.4N^3 + 25N^2 + 8N + 17$ statements.
 - Consider the runtime when N is *extremely large* .
 - We ignore constants like 25 because they are tiny next to N .
 - The highest-order term (N^3) dominates the overall runtime.
- We say that this algorithm runs "on the order of" N^3 .
- or **$O(N^3)$** for short ("Big-Oh of N cubed")

Complexity classes

- **complexity class:** A category of algorithm efficiency based on the algorithm's relationship to the input size N .

Class	Big-Oh	If you double N, ...	Example
constant	$O(1)$	unchanged	10ms
logarithmic	$O(\log_2 N)$	increases slightly	175ms
linear	$O(N)$	doubles	3.2 sec
log-linear	$O(N \log_2 N)$	slightly more than doubles	6 sec
quadratic	$O(N^2)$	quadruples	1 min 42 sec
cubic	$O(N^3)$	multiplies by 8	55 min
...
exponential	$O(2^N)$	multiplies drastically	$5 * 10^{61}$ years

Complexity classes



Collection efficiency

- Efficiency of our `ArrayIntList` or Java's `ArrayList`:

Method	<code>ArrayList</code>
<code>add</code>	$O(1)$
<code>add(index, value)</code>	$O(N)$
<code>get</code>	$O(1)$
<code>remove</code>	$O(N)$
<code>set</code>	$O(1)$
<code>size</code>	$O(1)$

Max subsequence sum

- Write a method `maxSum` to find the largest sum of any contiguous subsequence in an array of integers.
 - Easy for all positives: include the whole array.
 - What if there are negatives?

index	0	1	2	3	4	5	6	7	8
value	2	1	-4	10	15	-2	22	-8	5

Largest sum: $10 + 15 + -2 + 22 = 45$

- (Let's define the max to be 0 if the array is entirely negative.)
- Ideas for algorithms?

Algorithm 1 pseudocode

```
maxSum(a) :  
    max = 0.  
    for each starting index i:  
        for each ending index j:  
            sum = add the elements from a[i] to a[j].  
            if sum > max,  
                max = sum.  
  
    return max.
```

index	0	1	2	3	4	5	6	7	8
value	2	1	-4	10	15	-2	22	-8	5

Algorithm 1 code

- What complexity class is this algorithm?
 - **$O(N^3)$** . Takes a few seconds to process 2000 elements.

```
public static int maxSum1(int[] a) {
    int max = 0;
    for (int i = 0; i < a.length; i++) {
        for (int j = i; j < a.length; j++) {
            // sum = add the elements from a[i] to a[j].
            int sum = 0;
            for (int k = i; k <= j; k++) {
                sum += a[k];
            }
            if (sum > max) {
                max = sum;
            }
        }
    }
    return max;
}
```

Flaws in algorithm 1

- Observation: We are redundantly re-computing sums.
 - For example, we compute the sum between indexes 2 and 5:
 $a[2] + a[3] + a[4] + a[5]$
 - Next we compute the sum between indexes 2 and 6:
 $a[2] + a[3] + a[4] + a[5] + a[6]$
 - We already had computed the sum of 2-5, but we compute it again as part of the 2-6 computation.
 - Let's write an improved version that avoids this flaw.

index	0	1	2	3	4	5	6	7	8
value	2	1	-4	10	15	-2	22	-8	5

Algorithm 2 code

- What complexity class is this algorithm?
 - **$O(N^2)$** . Can process tens of thousands of elements per second.

```
public static int maxSum2(int[] a) {  
    int max = 0;  
    for (int i = 0; i < a.length; i++) {  
        int sum = 0;  
        for (int j = i; j < a.length; j++) {  
            sum += a[j];  
            if (sum > max) {  
                max = sum;  
            }  
        }  
    }  
    return max;  
}
```

index	0	1	2	3	4	5	6	7	8
value	2	1	-4	10	15	-2	22	-8	5

A clever solution

- *Claim 1* : A max range cannot start with a negative-sum range.

i	...	j	j+1	...	k
< 0			$\text{sum}(j+1, k)$		
$\text{sum}(i, k) < \text{sum}(j+1, k)$					

- *Claim 2* : If $\text{sum}(i, j-1) \geq 0$ and $\text{sum}(i, j) < 0$, any max range that ends at $j+1$ or higher cannot start at any of i through j .

i	...	j-1	j	j+1	...	k
≥ 0			< 0	$\text{sum}(j+1, k)$		
< 0				$\text{sum}(j+1, k)$		
$\text{sum}(?, k) < \text{sum}(j+1, k)$						

- Together, these observations lead to a very clever algorithm...

Algorithm 3 code

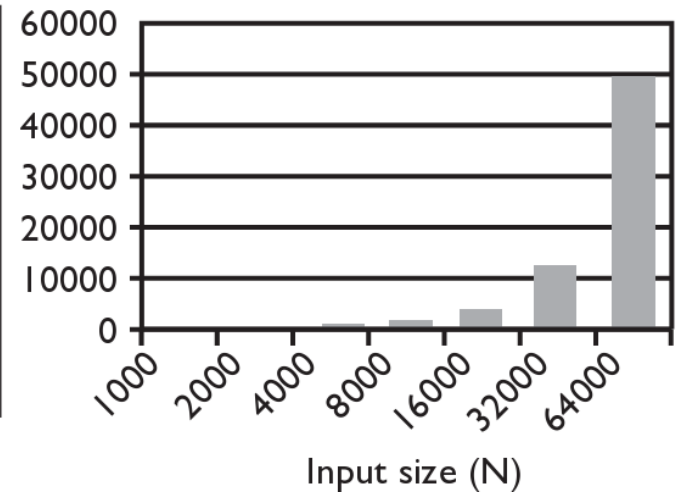
- What complexity class is this algorithm?
 - **O(N)**. Handles many millions of elements per second!

```
public static int maxSum3(int[] a) {
    int max = 0;
    int sum = 0;
    int i = 0;
    for (int j = 0; j < a.length; j++) {
        if (sum < 0) {    // if sum becomes negative, max range
            i = j;      // cannot start with any of i - j-1
            sum = 0;    // (Claim 2)
        }
        sum += a[j];
        if (sum > max) {
            max = sum;
        }
    }
    return max;
}
```

Runtime of first 2 versions

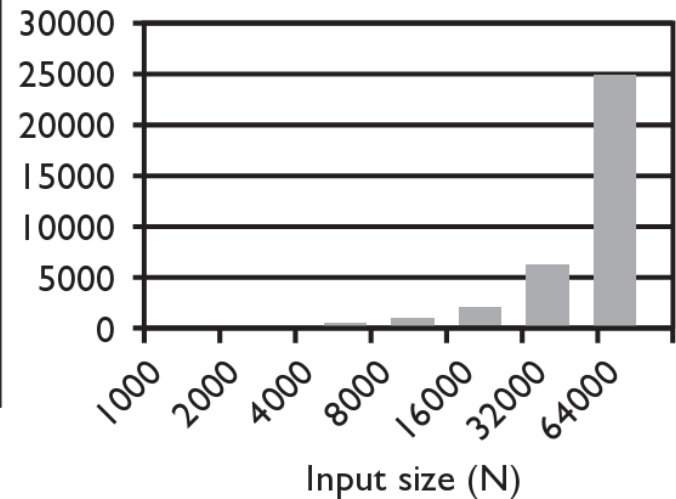
- Version 1:

N	Runtime (ms)
1000	15
2000	47
4000	203
8000	781
16000	3110
32000	12563
64000	49937



- Version 2:

N	Runtime (ms)
1000	16
2000	16
4000	110
8000	406
16000	1578
32000	6265
64000	25031



Runtime of 3rd version

- Version 3:

N	Runtime (ms)
1000	0
2000	0
4000	0
8000	0
16000	0
32000	0
64000	0
128000	0
256000	0
512000	0
1e6	0
2e6	16
4e6	31
8e6	47
1.67e7	94
3.3e7	188
6.5e7	453
1.3e8	797
2.6e8	1578

